

EMERGING COMMUNITIES IN NETWORKS — A FLOW OF TIES*

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Algorithms for search of communities in networks usually consist in discrete variations of links. Here we discuss a flow method, driven by a set of differential equations. Two examples are demonstrated in detail. First is a partition of a signed graph into two parts, where the proposed equations are interpreted in terms of removal of a cognitive dissonance by agents placed in the network nodes. There, the signs and values of links refer to positive or negative interpersonal relationships of different strength. Second is an application of a method akin to the previous one, dedicated to communities identification, to the Sierpiński triangle of finite size. During the time evolution, the related graphs are weighted; yet at the end, the discrete character of links is restored. In the case of the Sierpiński triangle, the method is supplemented by adding a small noise to the initial connectivity matrix. By breaking the symmetry of the network, this allows to a successful handling of overlapping nodes.

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1. Introduction

Reliable and fast methods of identification of communities in networks are of interest for numerous applications in potentially all branches of knowledge, from biology to computer sciences [1, 2]. For large networks, the condition of speed is crucial; on the other hand, to check all possible partitions is a NP-complete problem [2]. Further, the definition of communities as ‘groups of nodes within which connections are dense, and between which connections are sparser’ [1], although intuitively plausible remains fuzzy. A quantitative

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method to distinguish between possible partitions of a given network is to calculate the so-called modularity Q for all considered solutions [3]; the one with the largest modularity is the proper one. Yet, to check that a partition is the proper one, we have to find it at first.

The purpose of this work is to report an idea that a social system, driven by a designed time evolution, is going to find the optimal partition by one's own. We mean that the algorithm of solving the problem is equivalent to modelling the actual dynamics of the considered system. The method is to solve numerically a set of $N(N - 1)/2$ differential equations, where N is the number of nodes of the network. Each equation is devoted to one element of the connectivity matrix. Solving differential equations numerically is computationally costly, hence rarely used, with [4] as an exception. The advantage is that the method leads deterministically to the sought after solution, *i.e.* to the partition most close to the initial state.

Below, two variants of the proposed equations are reported. In the first variant, the matrix elements x_{ij} are related to social contacts between N persons. The intensity and character (friendly or hostile) of these contacts are given as absolute values and signs of x_{ij} . The obtained partition of the group is interpreted as the solution of the Heider balance problem. In a nutshell, the idea can be reported as follows. In a seminal paper [5], Heider has established the balanced and unbalanced configurations or mutual relationships of a triad of persons: in the former, the product of three related relationships is positive, and in the latter, it is negative. This work was generalized to a network by Harary [6], who has proved that the balanced state demands a clear division of the whole network into two groups, with friendly relationships within the groups and hostile relationships between all members of different groups. This concept has been supplemented by a dynamic aspect by Festinger [7], who indicated the ways people remove the cognitive dissonance, caused by unbalanced relationships. More recently, discrete algorithms of obtaining the balanced state have been constructed by [8, 9]. However, these prescriptions were plagued by jammed states, where some unbalanced triads were present. Later, our method of differential equations [10–12] was also investigated by Strogatz and coworkers [13], and no jammed states have been found there. In parallel, more and more historical events have been described in terms of the Heider balance [14, 15], with the formation of coalitions of European states before WWI as a canonical example.

In the second variant, a more general algorithm is described, which is appropriate to any number of communities — not just two. On the contrary to the previous one, all links are positive. An additional parameter β is introduced; its purpose is to separate relevant links from irrelevant ones. In a social system, the time evolution of links — relations between individuals — is equivalent to an improvement the relation in a dyad if its both members

have good relations with other persons, and to a deterioration of the relation in the dyad if the relations of both members with others are bad. The method is to be applied with different values of β , and the final result is selected as to obtain the largest value of the modularity index Q . Also, during the time evolution, the system passes through a series of subsequent partitions; again, the solution with largest Q is to be selected. The method has been formulated and validated for dense and sparse graphs in [16, 17]. A comparison with the results of [3] allowed to state that in many cases the new method prevails. Let us note that in both our variants, the time evolution of the network is entirely deterministic, with the actual data as an initial state.

In two subsequent sections, we describe examples of applications of the two variants of the equations. Both examples have been reported already in [11, 18]; here, we give a more detailed description of all stages of calculation. The first variant of the method has been validated in [11]. New element here is to determine the range of parameters where the solution is stable. Section 3 is devoted to the second variant, where it is applied to a particularly difficult case — a symmetric fractal. New element here is the time dependence of the modularity in the presence of noise. In the last section, we highlight the difficulty met in our case by the deterministic method, when the condition of symmetry is compared to an unstable fixed point of the system dynamics.

2. Cognitive dissonance

According to the concept of cognitive dissonance, persons involved in mutual social relationships tend to order them as to reach a consistent division of others into friends and enemies. This tendency, when expressed in terms of the symmetric connectivity matrix x_{ij} , means that the relationships tend to fulfil the condition $x_{ij}x_{jk}x_{ki} > 0$ for each triad i, j, k [5]. This means that either all relationships in a triad should be positive (friendly), or two of them should be negative (hostile) and one — positive. These two configurations are balanced, while the other two (one or three negative relationships) are unbalanced.

If a network configuration is unbalanced, what should be the rule to remove the dissonance? The idea is taken from the set of four simple statements: (a) *We like someone who likes someone that we like*; (b) *we like someone who dislikes someone we dislike*; (c) *we dislike someone who likes someone we dislike*; and (d) *we dislike someone who dislikes someone we like*. These statements are taken from the summary of the results of a careful laboratory experiment by Aronson [19]. The appropriate dynamics is that if A likes B and B likes D, then the relationship between A and D should

go better. If, on the contrary, A likes B and B dislikes D, the relationship between A and D gets worse. Perhaps the simplest rule of this kind is

$$\frac{dx_{ij}}{dt} = \sum_{k \neq i,j} x_{ik}x_{kj}. \quad (1)$$

The flaw of this equation is that some matrix elements go to infinity in a finite time. This is seen, for example, for a triad, if all relationships x, y, z are positive and initially equal; the solution is that $x(t) = x(0)/(1 - x(0)t)$, and so is $y(t)$ and $z(t)$. A peculiar case is when all relationships are equal and negative; then all links increase and all should become friends. Yet, because of the symmetry, all matrix elements decrease equally; in the effect, the relationships become just more and more neutral (*i.e.* close to zero) when time goes to infinity. Both these symmetrical, then non-generic, cases can be easily generalized to larger networks; just the time should be multiplied by the factor $(N - 2)$, where N is the number of nodes. Anyway, to evade the question how to interpret an infinite relationship, we limit the solutions x_{ij} to the range $(-R, R)$, where $R > 0$ is a parameter. This is done by the correction of the equation (1) as follows

$$\frac{dx_{ij}}{dt} = F(x_{ij}) \sum_{k \neq i,j} x_{ik}x_{kj}, \quad (2)$$

where $F(x) = 1 - (x/R)^2$. Actually, the analytical results [13] are obtained for the simplified version of equations, as in Eq. (1). In the non-generic symmetric case, when all relationships are initially of the same value, the marginally stable fixed point $x^* = 0$ of Eq. (1) is supplemented by two new fixed points for Eq. (2): unstable $x^* = -R$ and stable $x^* = +R$. In agreement with this, for $x^* = -R$ all triads are unbalanced, while for $x^* = +R$ the whole network is balanced and all relationships are friendly.

Both for Eqs. (1) and (2), once the network is balanced, it remains balanced forever; this is so, because in the balanced state $\text{sgn}(x_{ij}) = \text{sgn}(dx_{ij}/dt)$, for all links ij . This can be shown easily as follows: in each balanced triad, the product of three links is positive. Therefore, for a negative link, the product of the remaining two is negative; also, for a positive link, the product of the remaining two is also positive. These products contribute to the time derivative of the link, *q.e.d.* Below in this section, we refer to numerical results obtained with Eq. (2).

The case to be reported here is the set of data on relationships between 34 members of an unspecified karate club, collected by Zachary in 70s [20]. Shortly after this search was performed, a conflict appeared between the club members, and the group was divided into two. These data, available

on the Internet [21], are of common interest as a playground for all authors of algorithms designed to identify communities in networks. In particular, Newman applied two own algorithms there [3, 22]. In the former case, one member was assigned differently than in the reality; in the latter case, the obtained results exactly match the actual partition of the club members.

Actually, the data of Zachary are presented in [21] in two forms. In a more detailed version ('matrix C' in [20]), the graph is weighted: the values of the matrix elements — integers from 0 to 7 — reflect the relative strength/weakness of the relationships in the club. In a reduced version ('matrix E' in [20]), all non-zero values are substituted by 1s. We have applied Eq. (2) to both these forms. However, to get some links (relationships) negative, we reduced all matrix elements by the same value ε . For the matrix C, $\varepsilon = 2.5$, while for the matrix E, $\varepsilon = 0.5$. Both matrices are taken as the initial values of x_{ij} for $34 \times 33/2 = 561$ differential equations. In both cases, the parameter $R = 5$. The obtained time dependences of x_{ij} are shown in Fig. 1 (relationships weighted from 0 to 7, decreased by $\varepsilon = 2.5$) and Fig. 2 (relationships reduced to 0 or 1, decreased by $\varepsilon = 0.5$).

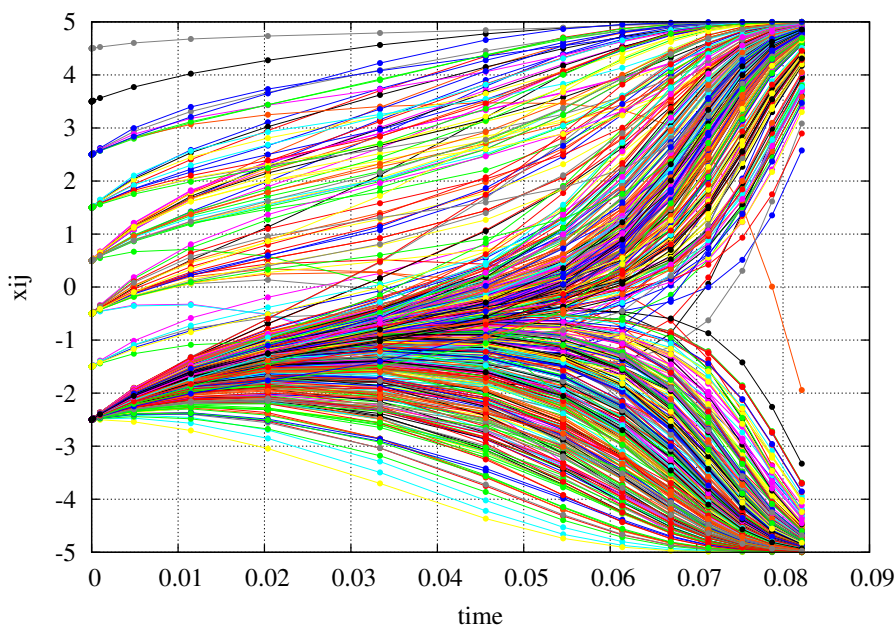


Fig. 1. Time dependence of the relationships x_{ij} driven by Eq. (2), with the elements of the matrix C (integers from 0 to 7) reduced by $\varepsilon = 2.5$ as the initial values. The parameter $R = 5.0$. When all relationships are balanced, the evolution is stopped.

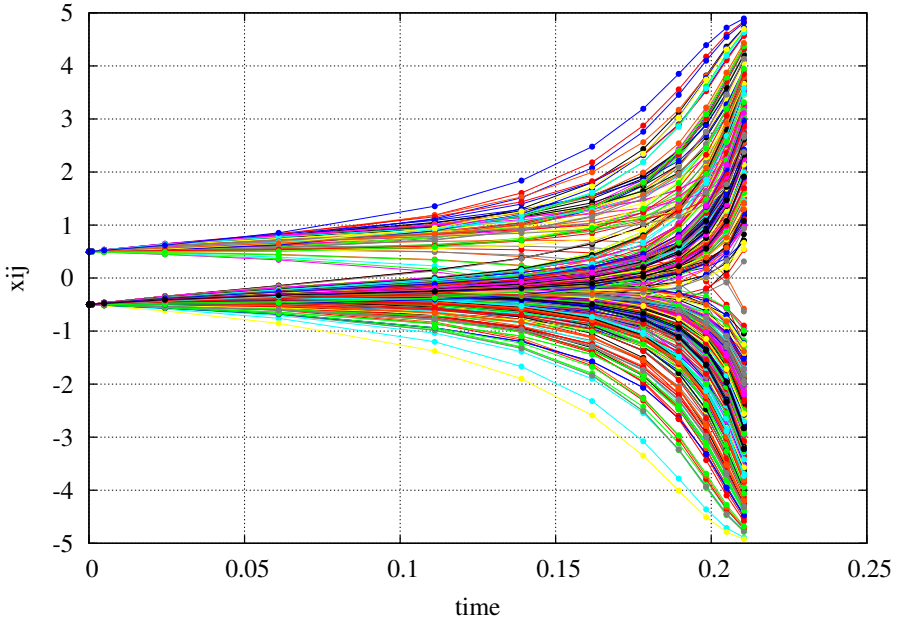


Fig. 2. Time dependence of the relationships x_{ij} driven by Eq. (2), with the elements of the matrix E (0 or 1) reduced by $\varepsilon = 0.5$ as the initial values. The parameter $R = 5.0$. When all relationships are balanced, the evolution is stopped.

The result is that Eq. (2) reproduces the actual partition of the members, except the member No. 9. (We note that in [11, 23], we erroneously wrote that our results exactly matched the real partition.) More surprisingly, our results for the matrices C and E are the same. This latter result is rather counterintuitive, because links equal to 1 or 2, classified as hostile in the matrix C, are converted to 1 in the matrix E, and then classified as friendly. Yet this means, that the solution is quite stable.

3. Noisy fractals

By definition, the Heider balance can be attained when either all relationships are friendly, or the network is divided into two parts. More generally, we need an equation which can lead to a division into any number of communities. The proposition [16, 17] is

$$\frac{dx_{ij}}{dt} = G(x_{ij}) \sum_{k \neq i, j} (x_{ik}x_{kj} - \beta), \quad (3)$$

where $G(x) = \Theta(x)\Theta(1-x)$ plays the same role as $F(x)$ in Eq. (2), and β is a parameter. Here, all links x_{ij} are positive; either they are weighted because a problem formulation, or they were originally zero or one, and belong to the intermediate range $(0, 1)$ only during the time evolution. The role of the parameter β is to separate meaningful values of the product $x_{ik}x_{kj}$ from negligible ones. If this product is larger than β , the pair of links (i, k) and k, j contributes to an increase of the link (i, j) ; if the product is smaller than β , this pair contributes to a reduction of x_{ij} . The value of the parameter has to be found by trials; the criterion is to get the modularity Q as large as possible. As the result of the time evolution, the modelled network is divided into more and more communities. The division which gives the largest Q is accepted as final.

Here, we are going to concentrate on an application of Eq. (3) to an identification of communities in a finite Sierpiński triangle. Namely, the network is formed from the nodes of the triangle, shown in Fig. 3. In [18], the search of this structure has been combined with the method of system compression [24, 25], applicable for symmetric systems. The Sierpiński triangle was a suitable example. In Fig. 4, we show an exemplary time dependence of the modularity. For better visibility, first 2880 time steps are not shown.

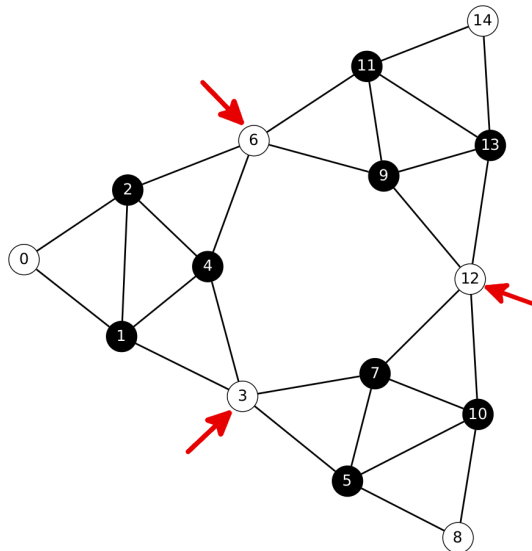


Fig. 3. The Sierpiński triangle of 15 nodes. Overlapping nodes are indicated by arrows.

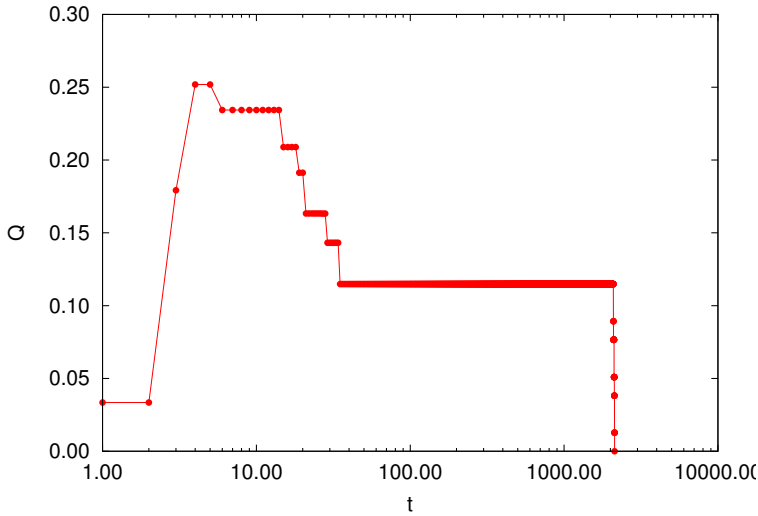


Fig. 4. Time dependence of the modularity Q — an example for the Sierpiński triangle of 15 nodes.

It is easy to indicate in Fig. 3 nodes which should belong simultaneously to two communities; yet, to be clear, we indicate these nodes by arrows. Any procedure designed to assign each node to one community must leave these nodes as communities of one node. Yet, this cannot be a criterion of overlapping, because more communities of one node are possible. Indeed, the application of Eq. (3) gives the structure as follows. There are three communities of three nodes (1,2,4), (5,7,10), (9,11,13) and six communities of one node (0), (3), (6), (8), (12), (14). How to distinguish the overlapping nodes?

The idea is to add some noise to the elements of the connectivity matrix of the investigated structure. With noise, the symmetry is broken and an overlapping node can be assigned to one or another community. Yet, the symmetry is preserved in the sense that the probabilities of assigning the node to two equivalent communities should be the same. Having performed the calculations many times, we can evaluate these probabilities. This is the criterion of overlapping nodes [18].

The noise is introduced by adding small random numbers ζ to the elements of the connectivity matrix of value zero, and by subtracting ζ from the elements of value one. The numbers ζ are different for each matrix element. Here, they are drawn from the range $(0, a)$ with uniform distribution; a is the amplitude of noise. Both a and β are free parameters, and the criterion to find their values is the maximization of the modularity Q . In Fig. 5, we show how the modularity Q , averaged over 10^4 trials, depends on

these parameters for $N = 15$. As can be deduced from the presented plots, $\beta = 0.15$ and the noise amplitude a of the order of 0.01 are appropriate. We note that the maximal Q exceeds its average value. To give an example, for the Sierpiński triangle of $N = 15$ nodes and the optimal values of parameters given above, the maximal Q is about 0.25. Yet, we must admit that this value is enhanced by some particular configuration of noise. More important result is that in the configurations which give this maximal Q , the nodes which are overlapping for $a = 0$ are reasonably included in the communities nearby. For the triangle from Fig. 4, a typical partition which gives large Q is (0,1,2,3,4), (5,7,8,10,12), (6,9,11,13,14). For this partition and zero noise, the modularity $Q = 0.26$.

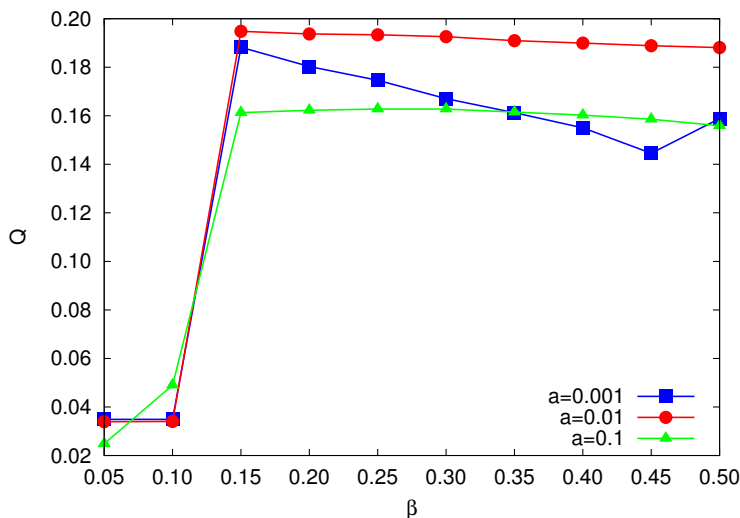


Fig. 5. The modularity Q as dependent on the parameter β for the Sierpiński triangle of 15 nodes, as in Fig. 3. Different curves are related to different amplitudes a of the applied noise. The results are averaged over 10^4 runs; for 10^3 runs, the picture is practically the same.

4. Discussion

To find a partition of a network is an inherently discrete task in the sense that all related variables are discrete. Both in the case of the Heider balance and in more general problem of communities, the task is to find a partition which is most close to the initial one. In the case of the Heider balance, the partition is such that positive links are only within the two parts of the network; these two parts are completely connected graphs with all links positive. All other links (between the parts) must be negative.

In the more general case, where all links are non-negative, the condition is that all links between different parts of network are zero. The idea to apply differential equations means that the transition from the initial to the final state is proceeding through a continuous variation of the links: they are represented by integers at the beginning and end of the process, and by reals in the meantime. The advantage is that the process is deterministic: there is no randomness included. As found by [13] for the case of the Heider balance, the dynamics driven by Eq. (2) always leads to the final state which obeys the demanded condition of balance, while an algorithm based on finite variations of discrete variables was shown to produce sometimes a jammed state. In terms of nonlinear dynamics, one could say that discrete algorithms based on Monte Carlo methods can produce solutions which are stable but unwanted, while differential equations are free from this flaw. The price we pay is the time of computation, much longer for differential equations than for discrete algorithms, as in [3, 22]. Also, while in the variant of the Heider balance the parameter ε is almost irrelevant, in the variant of more communities the simulation should be repeated for different values of the parameter β , between 0 and 1. The accepted value of β is the one which gives the largest modularity Q . Besides this, the algorithm is short and simple.

As we know from considerations of nonlinear dynamics, a system can be stuck in an unstable fixed point. We have seen a consequence of this in the case of the Heider balance, where the condition of symmetry (all matrix elements equal and negative) leads to a marginally stable fixed point where all links are zero. The condition of symmetry appears to be harmful also for the more general case. Namely, when we apply the method to a highly symmetric system, as the Sierpiński triangle, the overlapping nodes are artificially isolated. This is so, because — by symmetry — they can be assigned neither to one, nor to another of the neighbouring communities. Hence the role of noise, which breaks the symmetry and allows to obtain partitions with remarkably higher values of the modularity, than without noise.

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